

① Polynomial Representation of Convolutional Codes

$$u = \dots u_{-2} u_0 u_1 u_2 \dots$$

$$c = \dots c_{-2} c_0 c_1 c_2 \dots$$

$$u(D) = \dots u_{-1} D^{-1} + u_0 + u_1 D + u_2 D^2 + \dots$$

$$c(D) = \dots + c_{-2} D^{-2} + c_0 + c_1 D + c_2 D^2 + \dots$$

$$G = \begin{pmatrix} G_0 & G_1 & \dots & G_m \\ & G_0 & G_1 & \dots & G_m \\ & & G_0 & \dots & G_m \\ & & & \dots & \\ & & & & G_m \\ & & & & & \dots \\ & & & & & & \dots \end{pmatrix}$$

$$c = uG$$

$$c(D) = u(D) G(D)$$

$$G(D) = G_0 + G_1 D + \dots + G_m D^m$$

$$c_j(D) = \sum_{i=1}^k u_i(D) g_{ij}(D)$$

$g_{ij}(D) \rightarrow$ generator polynomial from i^{th} (input) to output (j^{th})

$$G(D) = \begin{pmatrix} g_{ij}(D) \end{pmatrix}$$

(2A)

$$G(D) = G_0 + G_1 D + \dots + G_m D^m$$

$$G(D) = (1 \ 2) + (1 \ 0) D + (1 \ 2) D^2 \\ = (1 + D + D^2 \quad 1 + D^2)$$

OR impulse responses are

$$h^{11} = (1 \ 2 \ 2) \quad h^{12} = (1 \ 2 \ 0 \ 2)$$

$$h^{11}(D) = 1 + D + D^2 \quad h^{12}(D) = 1 + D^2$$

$$g^{11}(D) = 1 + D + D^2 \quad g^{12}(D) = 1 + D^2$$

called generation polynomials

$$G(D) = (g_{ij}(D))$$

$$G(D) = (1 + D + D^2 \quad 1 + D^2)$$

2B

$$g^{11}(D) = 1 + D + D^2 \quad g^{12}(D) = 1 + D^2$$

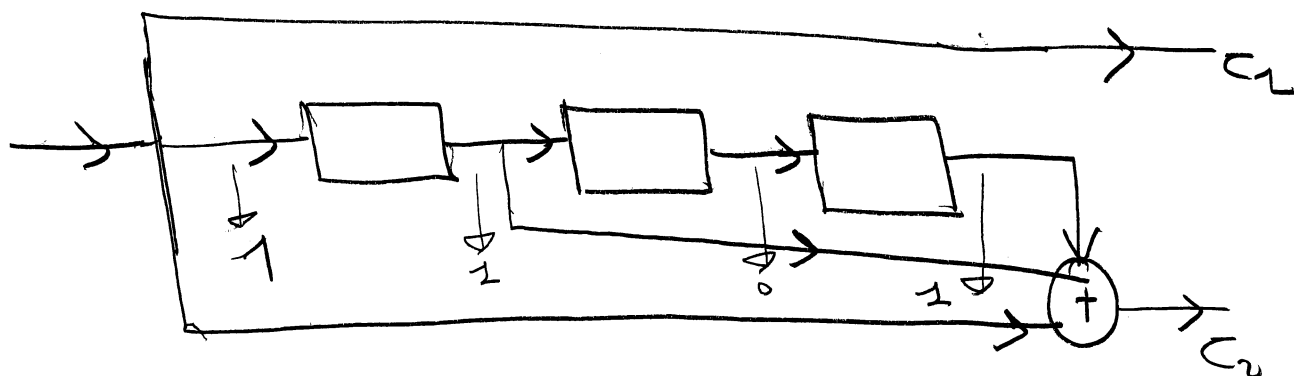
$$\sigma(D) = (1 + D + D^2 \quad 1 + D^2)$$

$$u(D) = (1 + D + D^2 + D^4)$$

$$c(D) = u(D) \sigma(D)$$

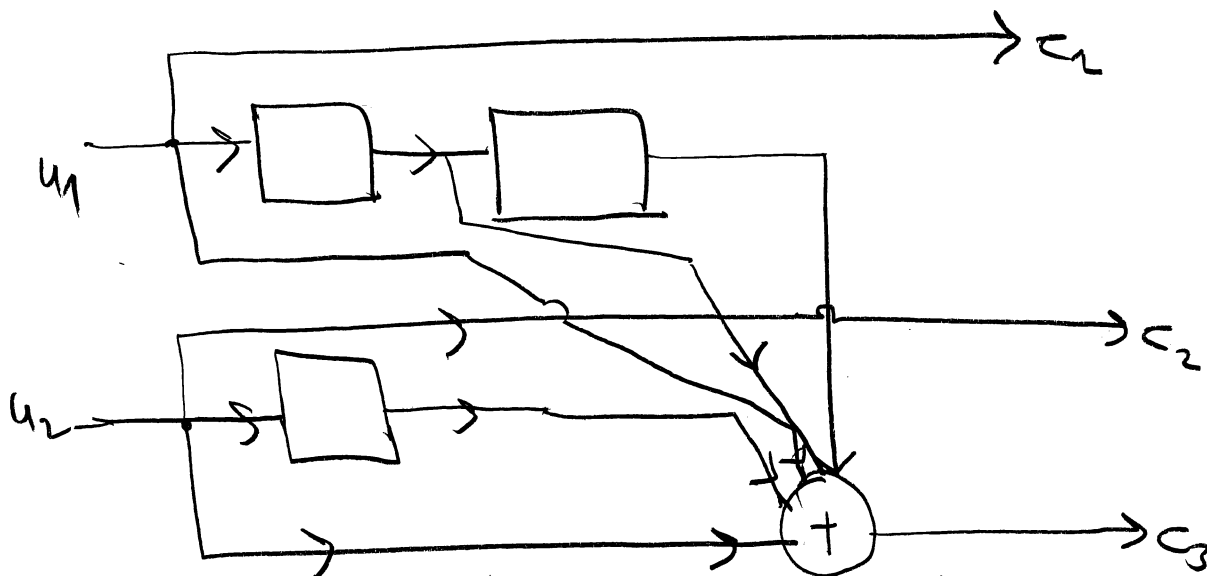
$$c_1(D) = u(D) g_{11}(D) \quad c_2(D) = u(D) g_{12}(D)$$

Ex 2



$$\sigma(D) = (1 \quad 1 + D + D^3)$$

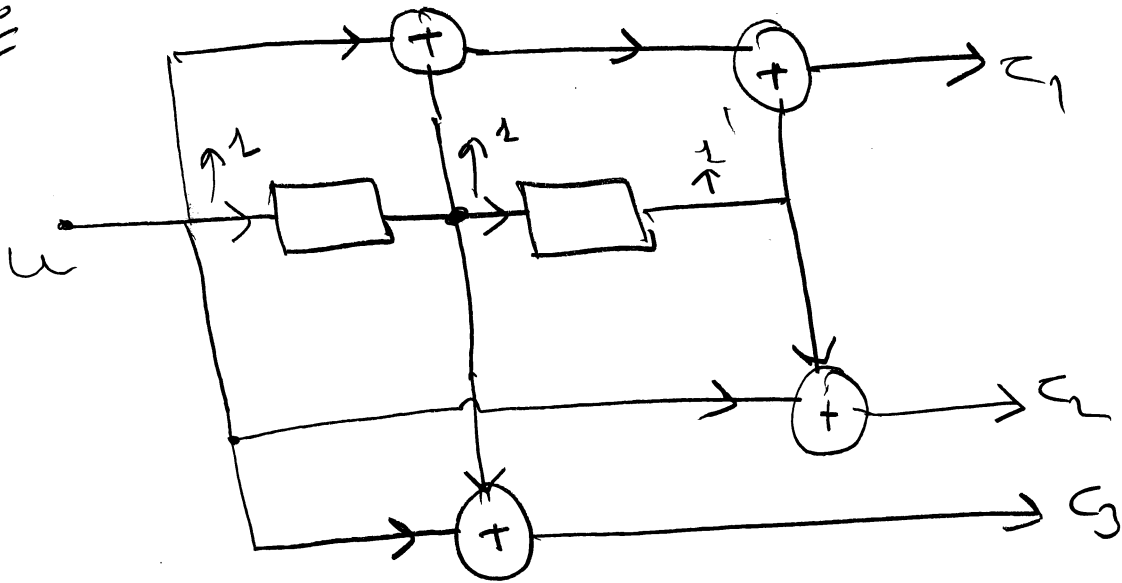
Ex 3



$$\sigma(D) = \begin{pmatrix} 1 & 0 & 1 + D + D^2 \\ 0 & 1 & 1 + D \end{pmatrix} \rightarrow H(D) = (1 + D^2 \quad 1 + D)$$

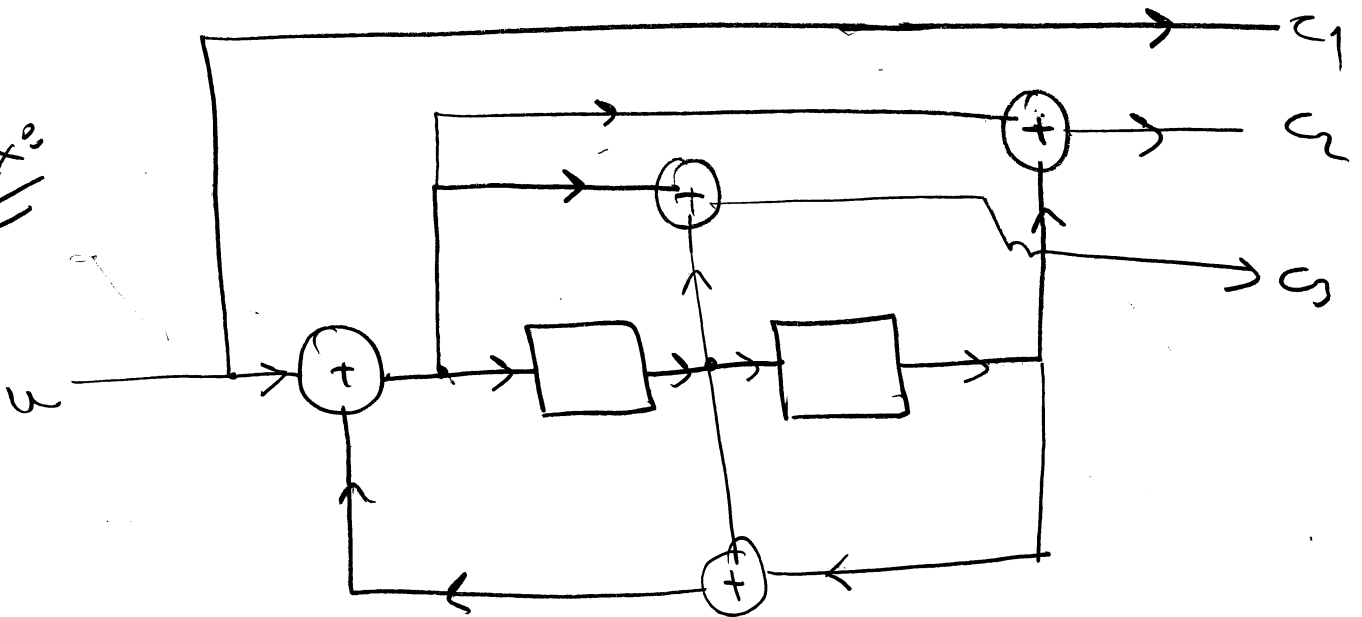
3

Exo



$$G(D) = \begin{bmatrix} 1+D+D^2 & 1+D^2 & 1 \end{bmatrix}$$

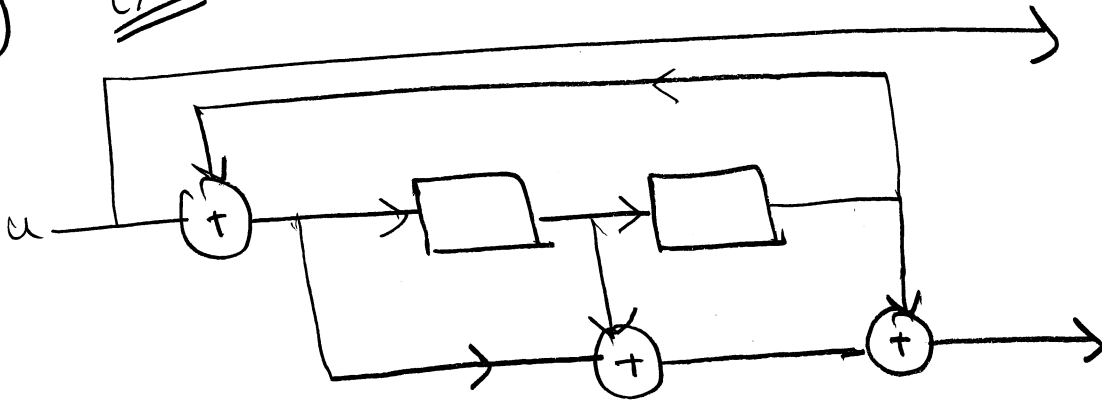
Exe



$$G(D) = \begin{bmatrix} 1 & \frac{1+D^2}{1+D+D^2} & \frac{1+D}{1+D+D^2} \end{bmatrix}$$

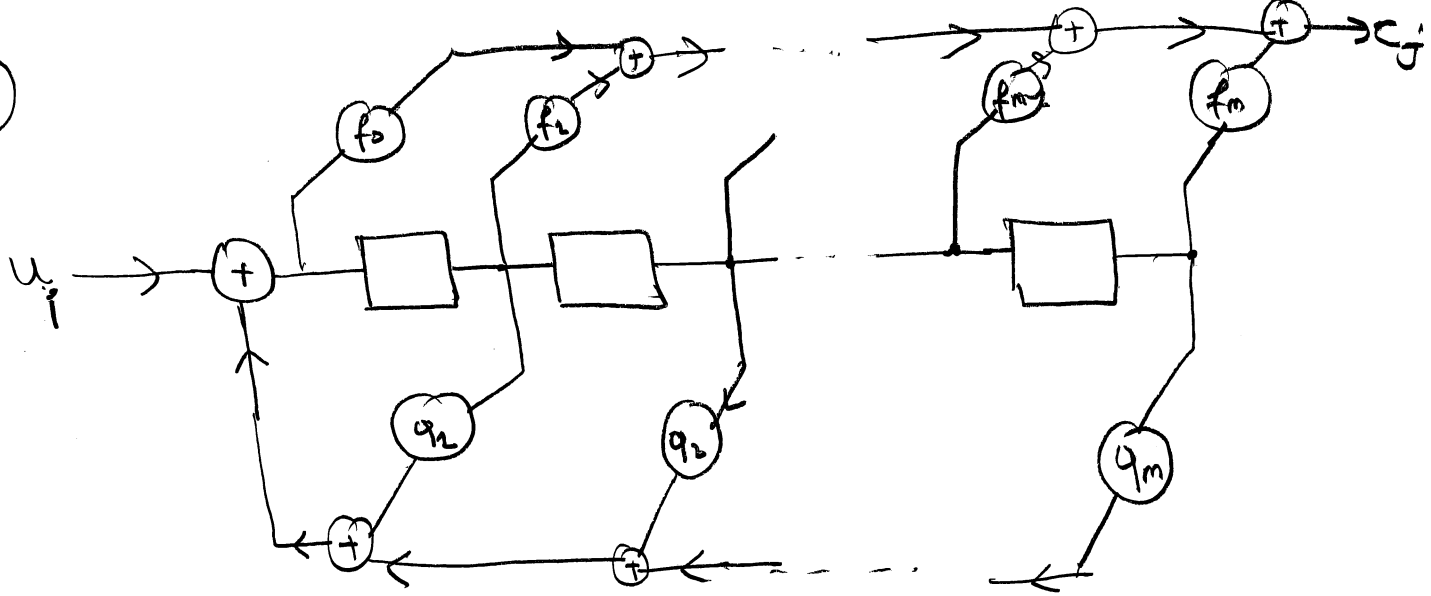
4

Ex:



$G(s) = ?$
Impulse response?

5

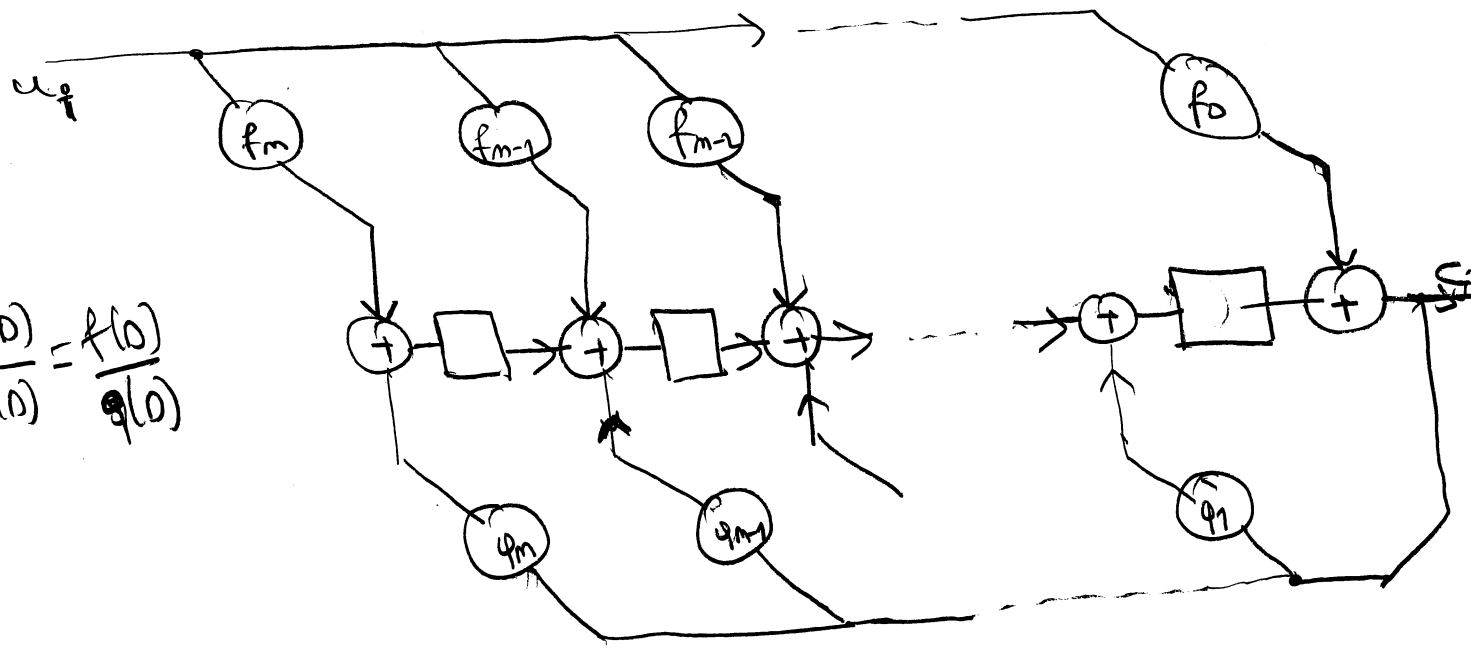


$$f(s) = f_0 + f_1 s + \dots + f_m s^m$$

$$q(s) = 1 + \varphi_1 s + \dots + \varphi_m s^m$$

↓
Controller
Canonical
Form

$$\frac{c_j(s)}{u_i(s)} = \frac{f(s)}{q(s)}$$



$$\frac{c_j(s)}{u_i(s)} = \frac{f(s)}{q(s)}$$

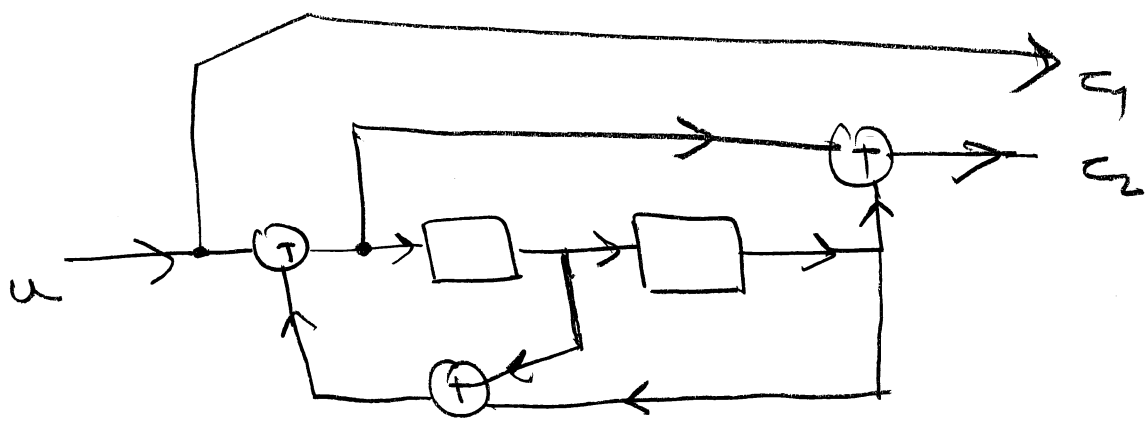
$$f(s) = f_0 + f_1 s + \dots + f_m s^m$$

$$q(s) = \varphi_2 s + \dots + \varphi_m s^m$$

↓ observer canonical
form

5A

Ex:



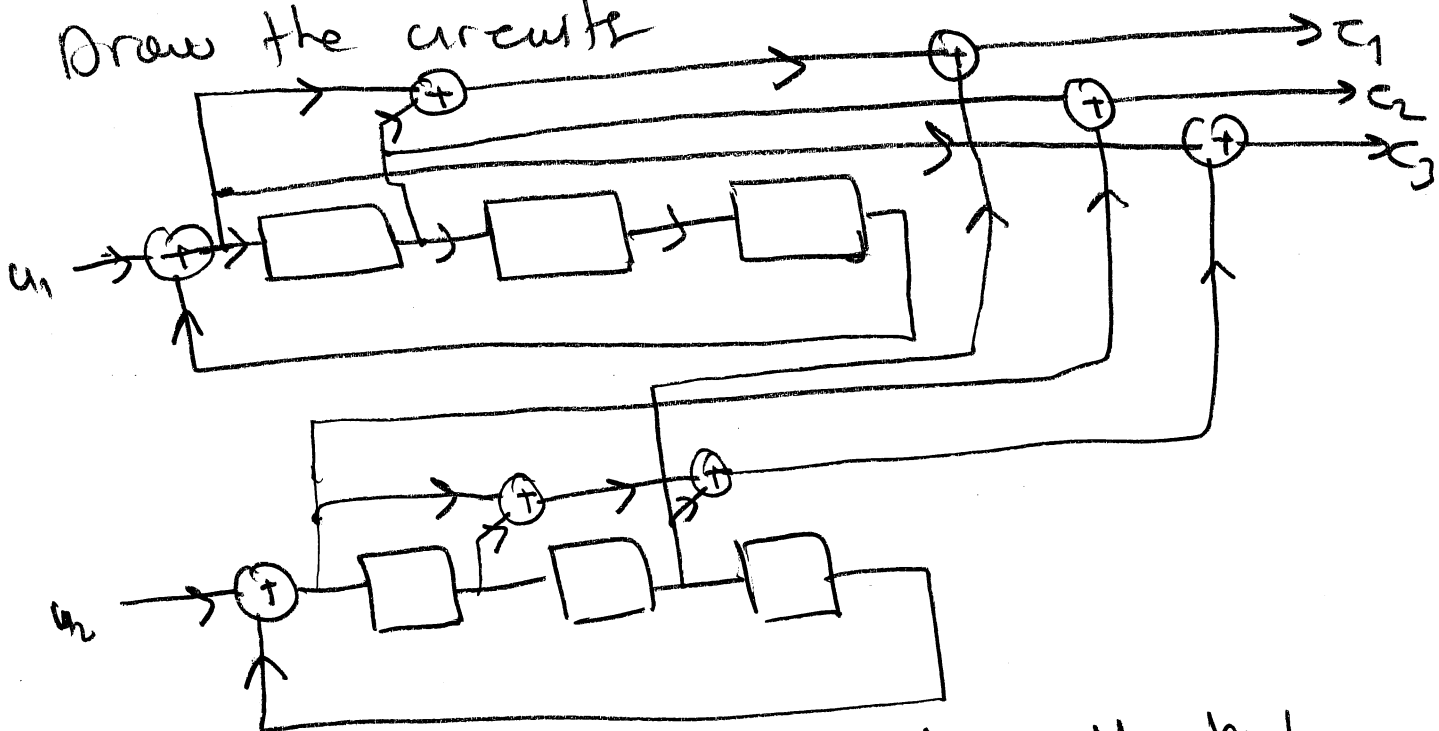
$$G(D) = \begin{pmatrix} 1 & \frac{1+D^2}{1+D+D^2} \end{pmatrix}$$

Ex:

$$G(D) = \begin{pmatrix} \frac{1}{1+D+D^2} & \frac{D}{1+D^3} & \frac{1}{1+D^3} \\ \frac{D^2}{1+D^3} & \frac{1}{1+D^3} & \frac{1}{1+D} \end{pmatrix}$$

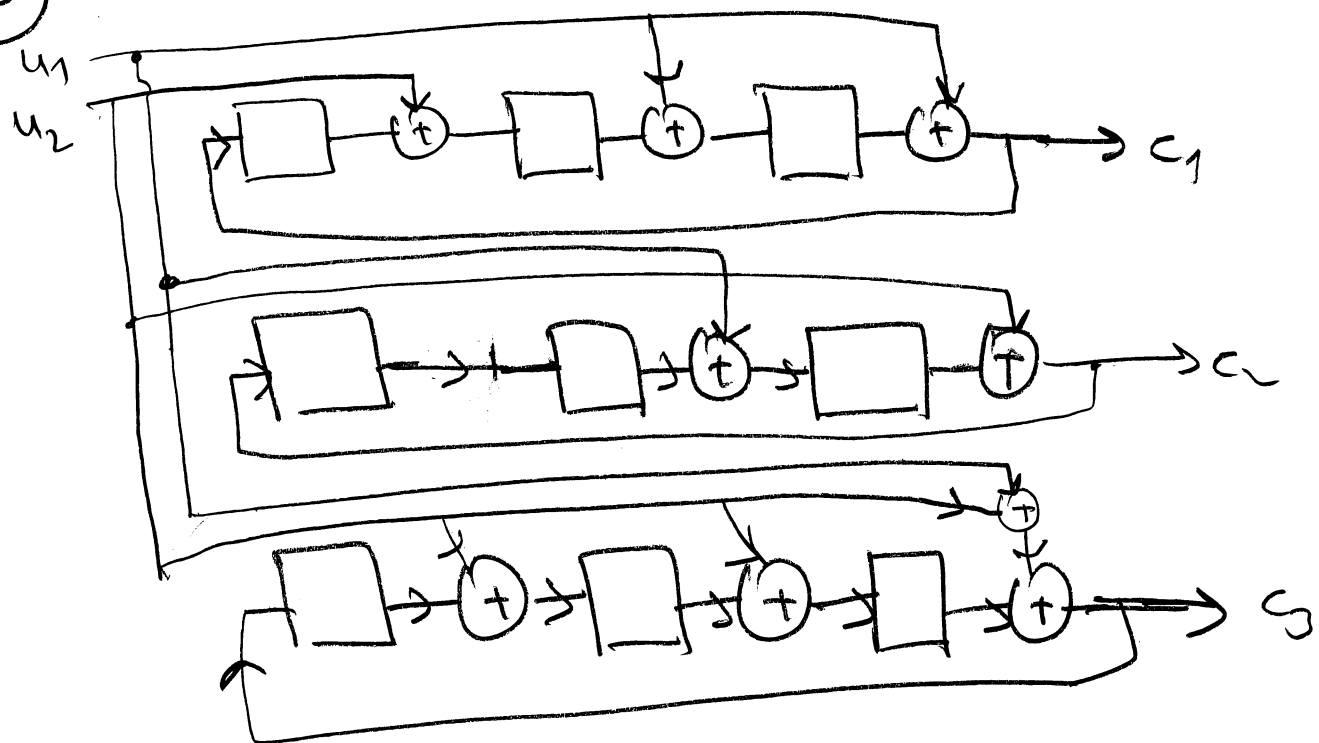
Now controller and observer canonical forms

Draw the circuits



↓ The solution has been taken from the text book "Fundamental of Convolutional Coding, Ziegengrov" Page: 37 But the solution seems to be wrong. Draw correct figures!

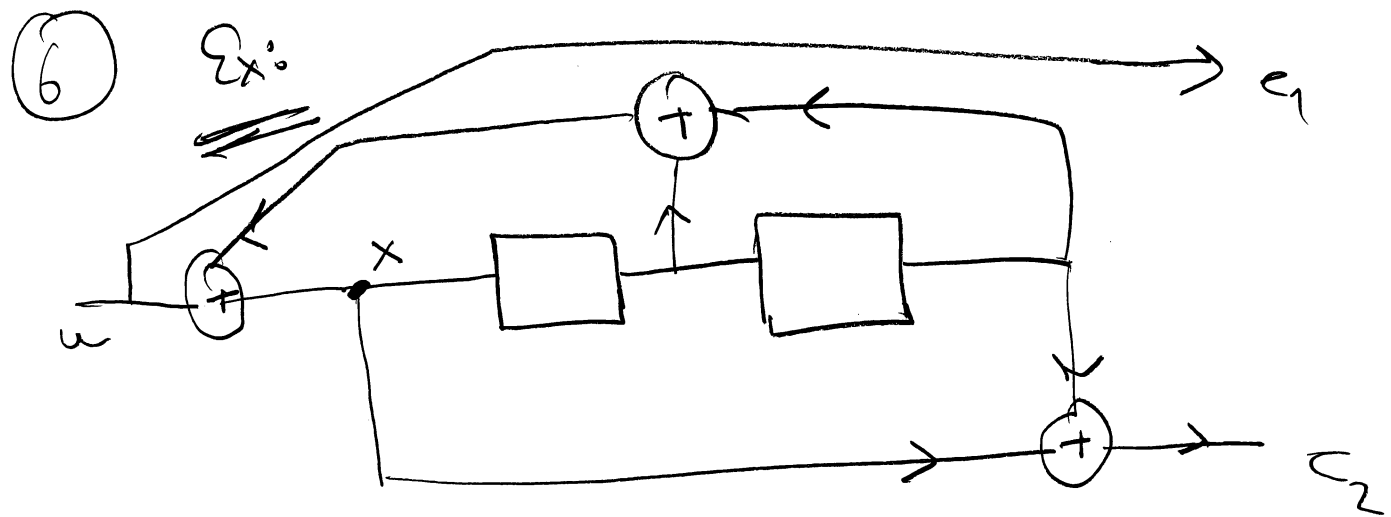
50



Ex 9

$$G(D) = \begin{pmatrix} 1+D & 0 & 1 \\ 0 & 1 & 1+D+D^2 \end{pmatrix}$$

Draw encoder circuits.



$$S_1 = 1 \quad S_2 = 1110110110\dots$$

$$x_t = u_t + x_{t-1} + x_{t-2} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} \rightarrow u_t = x_t + x_{t-1} + x_{t-2}$$

$$c_{2,t} = x_t + x_{t-2}$$

$$c_{1,t} = u_t$$

$$c_{2,t} = x_t + x_{t-2}$$

$$c_{2,t-1} = x_{t-1} + x_{t-3}$$

$$c_{2,t-2} = x_{t-2} + x_{t-4}$$

+

$$c_{2,t} + c_{2,t-1} + c_{2,t-2} = \underline{x_t + x_{t-1} + x_{t-2} + x_{t-4}}$$

$$u_t + u_{t-2}$$

7 Apply z-transform

$$\sum c_{2,t} D^t + \sum c_{2,t-1} D^t + \sum c_{2,t-2} D^t \\ = \sum u_t D^t + \sum u_{t-2} D^t$$


$$C_2(D) + D C_2(D) + D^2 C_2(D) = u(D) + D^2 u(D)$$

$$C_2(D) = \frac{1+D^2}{1+D+D^2} u(D)$$

$$\frac{C_2(D)}{u(D)} = \frac{1+D^2}{1+D+D^2}$$

The generator matrix is

$$G(D) = \left(1 \quad \frac{1+D^2}{1+D+D^2} \right) \rightarrow (1, 5/7)_{\text{octal}}$$

$$\left(1 \quad \frac{201}{111} \right)$$


8A

Hw

For $(1, 5/7)$ conv. code

$(1, 7/5)$ conv. code

a) Draw encoder circuit

b) Draw state diagram

State Table

c) Encode 11111 using
Trellis diagram, and show

the encoder path in trellis diagram

d) add trellis termination bits to 11111,

so that your encoder terminates

at 0 state

9

Catastrophic Convolutional Codes

- A catastrophic code maps information sequences with infinite Hamming weight to code sequences with finite Hamming weight.
- For a catastrophic code, a finite number of transmission errors can cause an infinite number of errors in the decoded information sequence. Hence, these codes should be avoided in practice.
- Systematic codes are never catastrophic.
- A loop in the state diagram that produces a zero output for a non-zero input is a necessary and sufficient condition for a catastrophic code.

Ex: (111) state has a loop associated with 1/00 in the following encoder

